

AN EMPIRICAL MODEL OF LOAD BALANCING TECHNIQUE IN SOFTWARE DEFINED NETWORK

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Abstract:

MULTIPATH LOAD BALANCE is an interesting research issue in the field of networks. Even though various models proposed by various authors. Every model has its own advantages and disadvantages. In this paper we are proposing efficient multi route path computation with in and out transmission between the nodes and Potential difference computes the time stamp difference between two nodes and number of hops between nodes. This approach generates the dynamic and secure path with less time complexity and ignores the unnecessary path

I.INTRODUCTION

SDN provides cloud-like processing inside a network. This provides network architects and directors to react rapidly to changes in business necessities through an incorporated control support that is disconnected from the actual structure of the network. As such, SDN makes a concentrated structure for the network that can impart and order the remainder of the network. SDN is utilized to make virtual overlay networks; programming characterized networks that depends on top of the actual structure framework.

There are three parts of the SDN networks, such as application, control and infrastructure. Communication between these layers and uses southbound and northbound APIs. In Northbound APIs, Applications using an SDN rely on the controller to tell them what the status of the network infrastructure is so that they can know what resources are available. Additionally, the SDN controller can automatically ensure application traffic is routed according to policies established by network administrators.

The applications talk to the control layer via the northbound APIs and tell the layer what resources the applications need, and their destination. The control layer orchestrates how the applications are given the resources available in the network. It also uses its intelligence to find the optimal path for the application in the context of its latency and security needs. Northbound APIs are often RESTful APIs. Orchestration is automated and not manually configured.

The SDN controller speaks with the network foundation, for example, switches a lot, through southward APIs. The network framework is determined what way the application information should take as chosen by the controller. Continuously, the controller can change how the switches a lot are moving information. The information no longer depends on the gadgets and directing tables to figure out where the information goes. All things being equal, the controller's insight settles on educated choices that advance the information's way.

A SDN controller is the product that gives a centralized perspective on and control over the whole network. Network administrators utilize the controller to oversee how the hidden framework's sending plane should deal with the traffic. The controller is likewise used to uphold arrangements that direct network conduct. Network administrators build up arrangements that are consistently applied to numerous hubs in the network.

Network approaches are deciding that are applied to traffic that figures out what level of access it has to the network, how much assets it is permitted, or what need it is relegated. Having a concentrated perspective on the network and the approaches set up makes for less difficult administration of the network that is more uniform and steadier.

There are benefits in SDN networks such as direct programmability, control over centralized network, agility and flexibility and reduced opex and capex.

II. RELATED WORK

Controller intended to build the readiness of the network by making it simple to oversee and adjust how traffic is taken care of. As a rule, the SDN Controller is the cerebrum of the SDN environment, imparting data down to the switches and switches with southward APIs, and up to the applications and business rationale with northward APIs. The Ryu Controller is upheld by NTT and is sent in NTT cloud server farms too.

The Ryu Controller provides software components, with well-defined application program interfaces (APIs), that make it easy for developers to create new network management and control applications. This component approach helps organizations customize deployments to meet their specific needs; developers can quickly and easily modify existing components or implement their own to ensure the underlying network can meet the changing demands of their applications.

The Ryu Controller source code is facilitated on GitHub and overseen and kept up by the open Ryu group. OpenStack, which runs an open cooperation zeroed in on building up a cloud

working framework that can control the register, storage, and networking resources of an association, underpins arrangements of Ryu as the Network Controller.

IT administrators compose explicit applications that convey to the Ryu controller on the most proficient method to communicate switches and switch. The Ryu Controller can utilize OpenFlow, or different conventions, to connect with the sending plane (switches and switches) to change how the network will deal with traffic streams. It has been tried and affirmed to work with various OpenFlow switches, including Open vSwitch and contributions from Centec, Hewlett Packard, IBM, and NEC.

III. DATA TRANSMISSION BETWEEN NODES

MULTIPATH:

Multipath routing is a routing technique that finds various courses to a destination in a network geography. By supporting numerous courses to a destination, network traffic can be disseminated genuinely through various ways in the network, or a component is known as burden adjusting, consequently expanding the proficiency of network utility. [4]

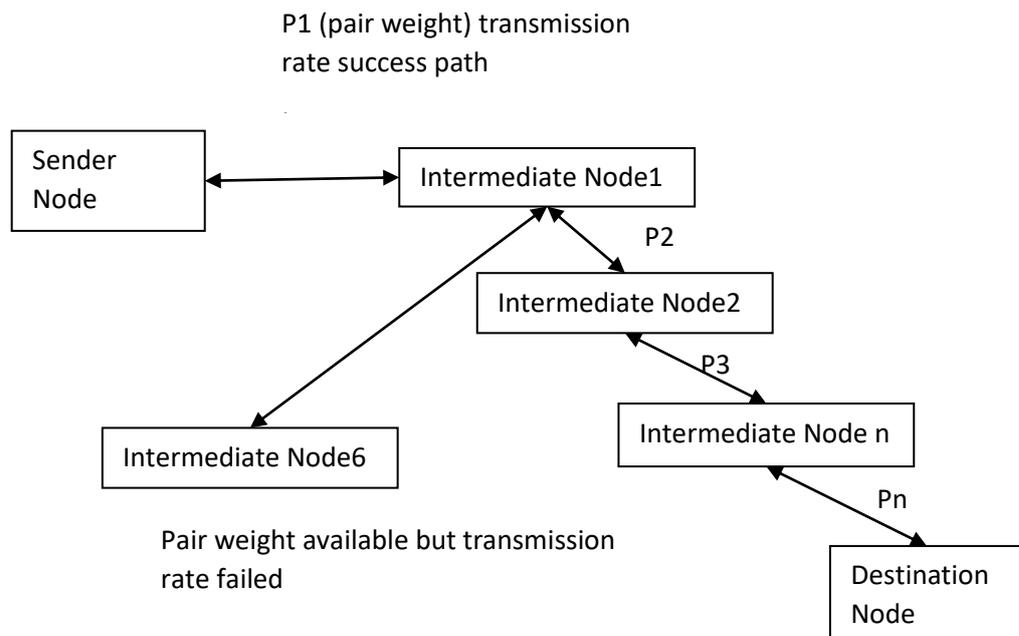


Fig 1 :Transmission between source to destination

Protocol of nodes:

This routing component improves the performance of the routing over TCP IP convention while transmission of information parcels from source to destination, by processing the ways from source to destination, different instrument utilizes different approach to speak with over network , each node contains its autonomous transmission in and out bundle subtleties.

Data transfer between in and out:

Potential contrast i.e time frame can be processed between transmission starting with one node then onto the next on the off chance that it is not exactly least limit disregard the node and pick another way, in any case check the following after node.

Info and yield Reliability bundle proportion table keeps up number of parcels moved however specific transfer node in last negligible stretch of time Rt_1 ,where R shows the hand-off node during the time interval t_1 ,here we consider the data bundles communicated in and out unwavering quality proportion and status 'S' ,it demonstrates hand-off node effectively sent the got parcels or not, for determination of ideal hand-off nodes in network.

Protocol of routing:

Nodes or vertices (V) and edges (E).the beneath chart shows $G(V,E)$, here v demonstrates the vertices and E shows the edge between the nodes shows the weight between source to destination node pair .C demonstrates absolute expense among source and destination node Identification if most brief path is a NP difficult issue, so we need to register paths of all conceivable arrangement until we arrive at briefest path. In our proposed model alongside cost factor we are thinking about the in and out dependability proportion of parcel (IOR) it should be fulfilled while calculation of cost of the nodes. IOR proportion can be kept up at a limit esteem, on the off chance that it meets least edge, it very well may be set to "Valid" ,so it tends to consider while cost calculation.

Allow use to think about a source node "A" needs to sends some data parcels to destination node "E" and B,C,D are moderate nodes, path can be founded on most elevated data rating by figuring normal of in out bundle transmission. The accompanying table shows test data rating table as follows.

In Byte data	Out Byte data
30	20
40	40
25	22
23	23

45	40
46	44

Data rating can be registered with normal of in and out concerning every single moderate node and data sent through most noteworthy rating path of certified nodes.

Routing Method:

Input: Start Node (SN), endNode(EN), Middle Node(MN), Cost=0, Path, I_t -In Transfer, O_t —Out transfer, avg_diff=0, threshold (t) -user defined

Output: Optimal cost, final_path

begin

Cost:=Cost + Avg_diff;

Path:=path+Path(SN,IN);

end

Next

Step3 : return moderate path to SN

Step1 : SN reads the in and out data values from next IN

Step2 :while (MN == EN)

GetIn_Out(MN)

Inavg := sum(I_t)/No of transactions;

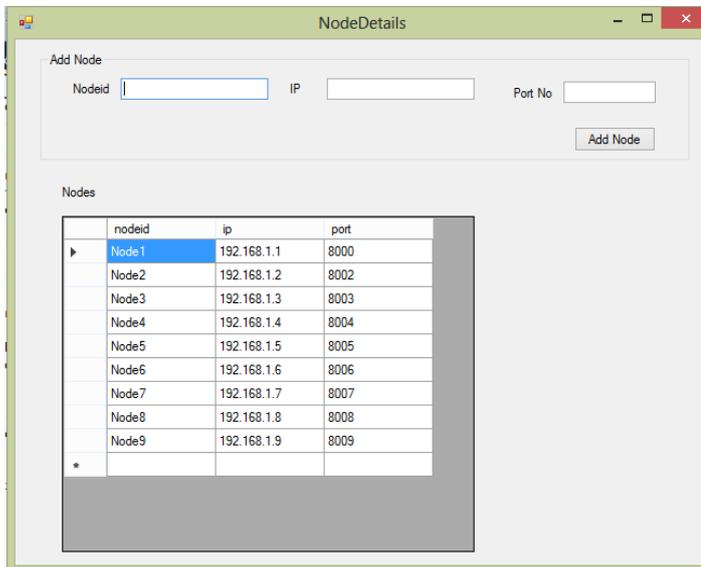
Outavg :=sum(O_t)/No of transactions;

Avg_diff:=Outavg-Inavg;

If(Avg_diff>threshold)

V.RESULTS

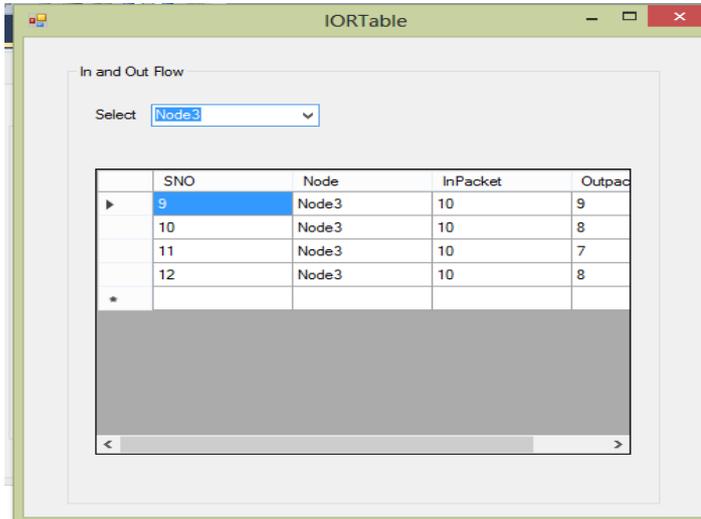
I) List of sample nodes which configured for simulation results.



nodeid	ip	port
Node1	192.168.1.1	8000
Node2	192.168.1.2	8002
Node3	192.168.1.3	8003
Node4	192.168.1.4	8004
Node5	192.168.1.5	8005
Node6	192.168.1.6	8006
Node7	192.168.1.7	8007
Node8	192.168.1.8	8008
Node9	192.168.1.9	8009

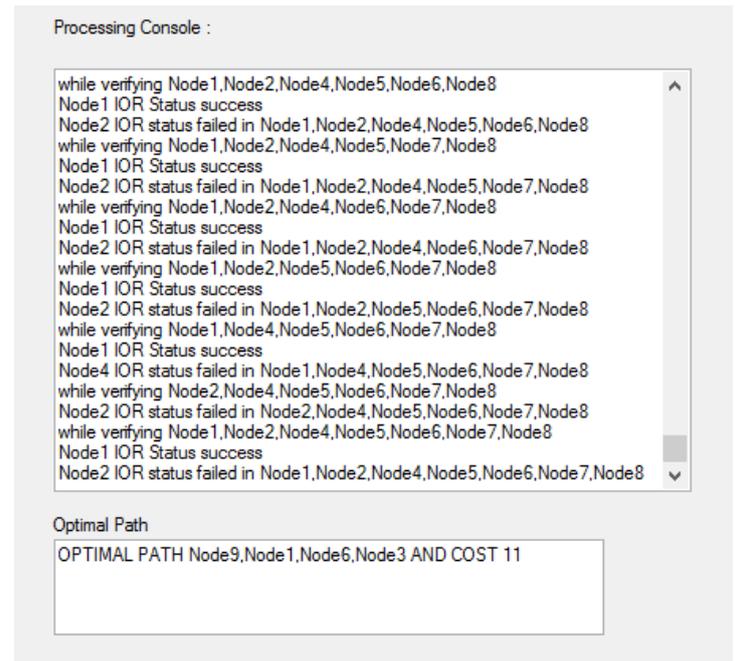
Fig2 : List of configured notes

II) In and out transmission rate configuration of node



SNO	Node	InPacket	Outpac
9	Node3	10	9
10	Node3	10	8
11	Node3	10	7
12	Node3	10	8

Fig3 : In and out transmission rate configuration of nodeIII) Communication cost as follows



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Processing Console :

while verifying Node1,Node2,Node4,Node5,Node6,Node8
Node1 IOR Status success
Node2 IOR status failed in Node1,Node2,Node4,Node5,Node6,Node8
while verifying Node1,Node2,Node4,Node5,Node7,Node8
Node1 IOR Status success
Node2 IOR status failed in Node1,Node2,Node4,Node5,Node7,Node8
while verifying Node1,Node2,Node4,Node6,Node7,Node8
Node1 IOR Status success
Node2 IOR status failed in Node1,Node2,Node4,Node6,Node7,Node8
while verifying Node1,Node2,Node5,Node6,Node7,Node8
Node1 IOR Status success
Node2 IOR status failed in Node1,Node2,Node5,Node6,Node7,Node8
while verifying Node1,Node4,Node5,Node6,Node7,Node8
Node1 IOR Status success
Node4 IOR status failed in Node1,Node4,Node5,Node6,Node7,Node8
while verifying Node2,Node4,Node5,Node6,Node7,Node8
Node2 IOR status failed in Node2,Node4,Node5,Node6,Node7,Node8
while verifying Node1,Node2,Node4,Node5,Node6,Node7,Node8
Node1 IOR Status success
Node2 IOR status failed in Node1,Node2,Node4,Node5,Node6,Node7,Node8
    
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Optimal Path
 OPTIMAL PATH Node9,Node1,Node6,Node3 AND COST 11

Fig 4: Computation of communication cost

V. CONCLUSION AND FUTURE SCOPE

In this paper we are addressing the multi route path issue with in and out transmission rate between in and out packet loads. In the previous paper, we have provided a clear idea of how to create a topology with five switches and two hosts experimental testbed along with its connectivity and flows analysis of keeping the Multiplemethod LoadBalancer performance as the central focus, but transmission between nodes should improved. That has been addressed in this paper with efficient cost communication mode. concluding our current research work with efficient over lay and underlay network, intermediate node can be chosen with transmission in and out rate average and data can be transmitted . Our Model gives more efficient results than traditional models.

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